

4.23 from BHM a)

Dual problem is

$$\begin{aligned}
 \min \quad & \sum_{i=1}^m b_i y_i + \sum_{j=1}^n u_j w_j \\
 \text{s.t.} \quad & \sum_i a_{ij} y_i + w_j \geq c_j \quad \forall j \\
 & y_i \text{ unrestricted } \forall i, \quad w_j \geq 0 \quad \forall j
 \end{aligned} \tag{18}$$

b) The optimality conditions are

$$\left. \begin{aligned}
 & \sum_{j=1}^n a_{ij} x_j = b_i \quad \forall i \\
 & x_j \leq u_j \quad \forall j \\
 & x_j \geq 0 \quad \forall j \\
 & \sum_{i=1}^m a_{ij} y_i + w_j \geq c_j \quad \forall j \\
 & w_j \geq 0 \quad \forall j, \\
 & y_i \text{ unrestricted } \forall i, \\
 & (u_j - x_j)w_j = 0 \quad \forall j \\
 & (c_j - (w_j + \sum_i a_{ij} y_i))x_j = 0 \quad \forall j
 \end{aligned} \right\} \begin{array}{l} \text{primal feasible} \\ \text{dual feasible} \\ \text{complementary slackness} \end{array} \tag{19}$$

c) Define $\bar{c}_j = c_j - \sum_i a_{ij}y_i$. Then the optimality conditions in part b imply

$$\begin{aligned}
 w_j &\geq \bar{c}_j \quad \forall j & (1) \\
 w_j &\geq 0 \quad \forall j & (2) \\
 (u_j - x_j)w_j &= 0 \quad \forall j & (3) \\
 (\bar{c}_j - w_j)x_j &= 0 \quad \forall j & (4)
 \end{aligned} \tag{20}$$

If $x_j = 0 \Rightarrow x_j < u_j \Rightarrow (3)$ implies $w_j = 0$ and (1) implies $\bar{c}_j \leq 0$.

If $x_j = u_j \Rightarrow x_j > 0 \Rightarrow (4)$ implies $w_j = \bar{c}_j$ and (2) implies $\bar{c}_j \geq 0$.

If $0 < x_j < u_j \Rightarrow (3)$ and (4) implies $w_j = 0$ and $\bar{c}_j = w_j$ and so $\bar{c}_j = 0$.

4.24 from BHM a)

Dual of the original LP is

$$\begin{aligned}
 \min \quad & \sum_{i=1}^m b_i y_i + \sum_{i=1}^q w_i d_i \\
 \text{s.t.} \quad & \sum_{i=1}^m y_i a_{ij} + \sum_{i=1}^q w_i e_{ij} \geq c_j \quad \forall j \\
 & y_i \geq 0, w_i \geq 0
 \end{aligned} \tag{21}$$

The optimality conditions for the original LP are

$$\begin{aligned}
 & \left. \begin{aligned}
 \sum_{j=1}^n a_{ij}x_j &\leq b_i \quad i = 1, \dots, m \\
 \sum_{j=1}^n e_{ij}x_j &\leq d_i \quad i = 1, \dots, q \\
 x_j &\geq 0 \quad \forall j
 \end{aligned} \right\} \text{primal feasible} \\
 \text{(Cond 1)} \quad & \left. \begin{aligned}
 \sum_{i=1}^m y_i a_{ij} + \sum_{i=1}^q w_i e_{ij} &\geq c_j \quad \forall j \\
 y_i &\geq 0, w_i \geq 0
 \end{aligned} \right\} \text{dual feasible} \\
 & \left. \begin{aligned}
 (-\sum_{j=1}^n a_{ij}x_j + b_i)y_i &= 0 \quad i = 1, \dots, m \\
 (-\sum_{j=1}^n e_{ij}x_j + d_i)w_i &= 0 \quad i = 1, \dots, q \\
 (-\sum_{i=1}^m y_i a_{ij} - \sum_{i=1}^q w_i e_{ij} + c_j)x_j &= 0 \quad \forall j
 \end{aligned} \right\} \text{complementary slackness}
 \end{aligned} \tag{22}$$

Next consider the Lagrangean problem. Define $\bar{c}_j = c_j - \sum_i a_{ij}y_i$. The dual of this problem is

$$\begin{aligned}
 \min \quad & \sum_{i=1}^q w_i d_i \\
 \text{s.t.} \quad & \sum_{i=1}^q w_i e_{ij} \geq \bar{c}_j \quad \forall j \\
 & w_i \geq 0
 \end{aligned} \tag{23}$$

The optimality conditions for the Lagrangean problem are

$$\begin{array}{l}
\left. \begin{array}{l} \sum_{j=1}^n e_{ij}x_j \leq d_i \quad i = 1, \dots, q \\ x_j \geq 0 \quad \forall j \end{array} \right\} \text{primal feasible} \\
\left. \begin{array}{l} \sum_{i=1}^q w_i e_{ij} \geq \bar{c}_j \quad \forall j \\ w_i \geq 0 \end{array} \right\} \text{dual feasible} \\
\left. \begin{array}{l} (-\sum_{j=1}^n e_{ij}x_j + d_i)w_i = 0 \quad i = 1, \dots, q \\ (-\sum_{i=1}^q w_i e_{ij} + \bar{c}_j)x_j = 0 \quad \forall j \end{array} \right\} \text{complementary slackness}
\end{array} \tag{24}$$

It is easy to see that if x_j^*, y_i, w_i satisfy Cond 1, then they also satisfy Cond 2, which implies that x_j^* also solves the Lagrangean problem.

4.30 from BHM Firstly, let's solve

$$\min_{y_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m) \tag{25}$$

and

$$\max_{x_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m). \tag{26}$$

We get

$$\min_{y_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m) = \begin{cases} \sum_{j=1}^n c_j x_j, & \text{if } -\sum_{j=1}^n a_{ij}x_j + b_i \geq 0, \forall i, \\ -\infty, & \text{otherwise,} \end{cases} \tag{27}$$

and

$$\max_{x_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m) = \begin{cases} \sum_{i=1}^m b_i y_i, & \text{if } c_j - \sum_{i=1}^m a_{ij}y_i \leq 0, \forall j \\ +\infty, & \text{otherwise.} \end{cases} \tag{28}$$

So

$$\max_{x_j \geq 0} \min_{y_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m) = \min_{y_j \geq 0} \max_{x_j \geq 0} f(x_1, \dots, x_n; y_1, \dots, y_m) \tag{29}$$

implies that the following two problems have the same optimal objective function value:

$$\begin{array}{l}
\max \quad \sum_{j=1}^n c_j x_j \\
s.t. \quad -\sum_{j=1}^n a_{ij}x_j + b_i \geq 0, \forall i \\
\quad \quad x_j \geq 0, \forall j,
\end{array} \tag{30}$$

$$\begin{aligned} \min \quad & \sum_{i=1}^m b_i y_i \\ \text{s.t.} \quad & c_j - \sum_{i=1}^m a_{ij} y_i \leq 0, \forall j \\ & y_i \geq 0, \forall j. \end{aligned} \tag{31}$$

However, note that (31) is just the dual problem to (30), so they have same optimal function value just implies the strong duality property.