

Lec 11: Strong duality Thm; Sensitivity

Note Title

10/14/2005

$$\text{Max } 2x_1 + 3x_2$$

$$-x_1 + x_2 \leq 5$$

$$x_1 + 3x_2 \leq 35$$

$$x_1 \leq 20$$

$$x_1, x_2 \geq 0$$

$$x_1 = 20 - s_3$$

$$x_2 = 5 - \frac{1}{3}s_2 - \frac{1}{3}s_3$$

$$s_1 = 20 + \frac{1}{3}s_2 - \frac{4}{3}s_3$$

$$Z = 55 - s_2 - s_3$$

$$\text{Min } 5y_1 + 35y_2 + 20y_3$$

$$-y_1 + y_2 + y_3 \geq 2$$

$$y_1 + 3y_2 \geq 3$$

$$y_1, y_2, y_3 \geq 0$$

$$y_2 = 1 - \frac{1}{3}y_1 + \frac{1}{3}w_2$$

$$y_3 = 1 + \frac{4}{3}y_1 + w_1 - \frac{1}{3}w_2$$

$$Z = -55 - 20y_1 - \frac{20w_1}{5} - 5w_2$$

Primal optimal soln.:

$$x_1^* = 20, x_2^* = 5$$

From last row of opt. dual dictionary

$$\text{Coeff. of } w_1 = -20$$

$$w_2 = -5$$

Dual optimal soln:

$$y_1^* = 0, y_2^* = 1, y_3^* = 1$$

from last row of opt. primal dictionary:

$$\text{Coeff. of } s_1 = 0$$

$$s_2 = -1$$

$$s_3 = -1$$

Observation:

Opt. Soln. to one problem =

negative of the coefficients of
Slack/Surplus variables in the last
row of the optimal dictionary
of the other problem.

Why?

$$Z = 55 - s_2 - s_3 - \otimes$$

We suspect $\{y_2 = 1, y_3 = 1, y_1 = 0\}$
is an optimal dual soln.

So:

$$Z = 55 - y_2 s_2 - y_3 s_3$$

BUT: $Z = 2x_1 + 3x_2$

$$s_2 = 35 - x_1 - 3x_2$$

$$s_3 = 20 - x_1$$

So:

$$2x_1 + 3x_2 = 55 - y_2(35 - x_1 - 3x_2) - y_3(20 - x_1)$$

As this holds for all x_1, x_2 , the coefficients of x_1, x_2 should be the same. & the "Constant" term should be zero. So:

$$\left. \begin{aligned} 2 &= y_2 + y_3 \\ 3 &= 3y_2 \end{aligned} \right\} \text{dual feasibility}$$

$$55 = 35y_2 + 20y_3 \leftarrow \text{Optimal dual soln. [Why?]}$$

$$\rightarrow y_1 = 0, y_2 = 1, y_3 = 1 \text{ is dual feasible}$$

$$\rightarrow \text{Obj value} = 55, \text{ hence dual optimal [Why?]}$$

In general:

$$\text{Max } \sum_j c_j x_j$$

$$\left[\begin{array}{l} i = 1, 2, \dots, m \\ j = 1, 2, \dots, n \end{array} \right]$$

$$\sum_j a_{ij} x_j \leq b_i$$

$$x_j \geq 0$$

Last row of opt. Primal dictionary:

$$Z = Z^* - \sum_j \bar{c}_j x_j - \sum_i \bar{d}_i S_i$$

By optimality of Primal dictionary

$$\bar{c}_j \geq 0, \quad \bar{d}_i \geq 0$$

for all j

for all i

Rewriting $\textcircled{*}$ using

$$S_i = b_i - \sum_j a_{ij} x_j$$

$$Z = \sum_j c_j x_j$$

$$\sum_j c_j x_j = Z^* - \sum_j \bar{c}_j x_j - \sum_i \bar{d}_i \left(b_i - \sum_j a_{ij} x_j \right)$$

$$= Z^* - \sum_j x_j \left\{ \bar{c}_j - \sum_i \bar{d}_i a_{ij} \right\}$$

$$- \sum_i b_i \bar{d}_i$$

Since this is valid for all

(x_1, x_2, \dots, x_n) , we must

have:

$$c_j = -\bar{c}_j + \sum_i \bar{d}_i a_{ij} \quad - \textcircled{1}$$

$$Z^* = \sum b_i \bar{d}_i \quad - \textcircled{2}$$

$$\textcircled{1} \Rightarrow \sum_i a_{ij} \bar{d}_i = c_j + \bar{c}_j \geq c_j$$

[because $\bar{c}_j \geq 0$]

Moreover, $\bar{d}_i \geq 0$.

So: \bar{d}_i is dual feasible

By $\textcircled{2}$, Obj. of dual at this feasible soln. $\{\bar{d}_i\} = Z^*$

So: $\{\bar{d}_i\}$ is dual optimal

Summary:

We just proved Strong duality
thm. for linear Programming:

If an LP has a finite opt-
value, then so does its dual
& the opt. values of the 2
problems are the same.

$$\text{Max } 60x_1 + 30x_2 + 20x_3$$

$$8x_1 + 6x_2 + x_3 + s_1 = 48$$

$$4x_1 + 2x_2 + 1.5x_3 + s_2 = 20$$

$$2x_1 + 1.5x_2 + \frac{1}{2}x_3 + s_3 = 8$$

$$x_1, x_2, x_3, s_1, s_2, s_3 \geq 0$$

$$s_1 = 24 + 2x_2 - 2s_2 + 8s_3$$

$$x_3 = 8 + 2x_2 - 2s_2 + 4s_3$$

$$x_1 = 2 - 1.25x_2 + \frac{1}{2}s_2 - 1.5s_3$$

$$Z = 280 - 5x_2 - 10s_2 - 10s_3$$

OPTIMAL DICTIONARY

$$x_B = B^{-1}b - B^{-1}N x_N$$

$$z = c_B B^{-1}b - (c_B B^{-1}N - c_N) x_N$$

Feasibility: $B^{-1}b \geq 0$

optimality: $[c_B B^{-1}N - c_N] \leq 0$

Ex: $B = \begin{bmatrix} 1 & 1 & 8 \\ 0 & 1.5 & 4 \\ 0 & 0.5 & 2 \end{bmatrix}$

$$B^{-1} = \begin{bmatrix} 1 & 2 & -8 \\ 0 & 2 & -4 \\ 0 & -\frac{1}{2} & -\frac{3}{2} \end{bmatrix}$$

Changing obj fn. coefficient

• $C_2 = 30$. Suppose $C_2 \uparrow$

ie. $C_2 = 30 + \delta$

1) For what values of δ is the current basis optimal?

[Range]

2) Suppose C_2 is changed beyond range. How to reoptimize?

B^{-1} , b remain unchanged

$x_2 = \text{non-basic} \Rightarrow C_B$ remains unchanged.

Only coefficient of x_2 in last row changes.

$$C_B B^{-1} = [0 \ 10 \ 10]$$

$$A_2 = x_2 \text{ column in } A = \begin{bmatrix} 6 \\ 2 \\ 1.5 \end{bmatrix}$$

$$\begin{aligned} \bar{c}_2 &= C_B B^{-1} A_2 - c_2 = [0 \ 10 \ 10] \begin{bmatrix} 6 \\ 2 \\ 1.5 \end{bmatrix} - (30 + \delta) \\ &= 35 - (30 + \delta) = 5 - \delta \end{aligned}$$

Thus:

* C_2 Can be decreased by any amount, and

* C_2 Can be increased by up to 5

without affecting the optimality of the current basis.

Also, the current soln. does not change.

∴ C_2 is increased beyond 5, current basis is not optimal.

Use Primal Simplex to
Continue.....